Lecture 6: Internet Layer

Reading 5.1. and 5.6 in Computer Networks, Tanenbaum

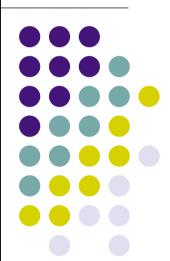




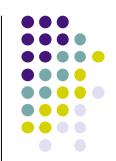
- Internet Protocol
- IP address and IP packet format
- ICMP- Protocol for control message

Introduction about IP

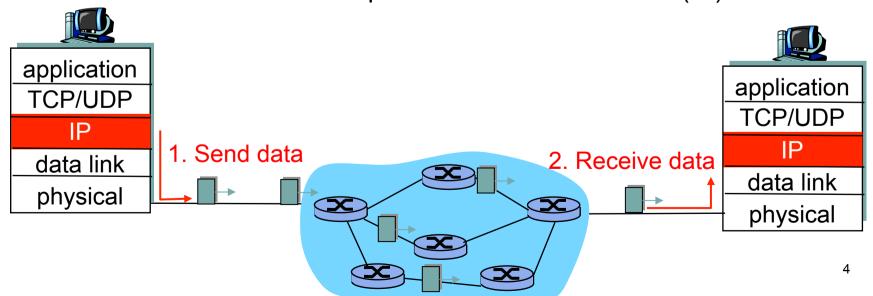
Concepts
Store and forward principles
Characteristic of IP



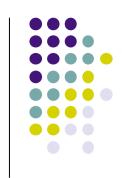
Network layer and Internet protocol



- Role of network layer: Transferring data between distant nodes
- Two main functionalities of Network layer
 - Routing: Determine the path for transferring data from the source to the destination nodes → Role of routing protocol.
 - Forwarding: Transferring data from the an incoming port to an outgoing port of a node (router) according to the path defined above → Role of routed protocol: Internet Protocol (IP)



Network layer and Internet protocol



- Layer 2 devices allow to connect limited number of close hosts
- When hosts are far from each other, intermediates nodes with forwarding and path finding functionality is needed → Router
 - Finding routes
 - Forwarding data according to destination Network layer address

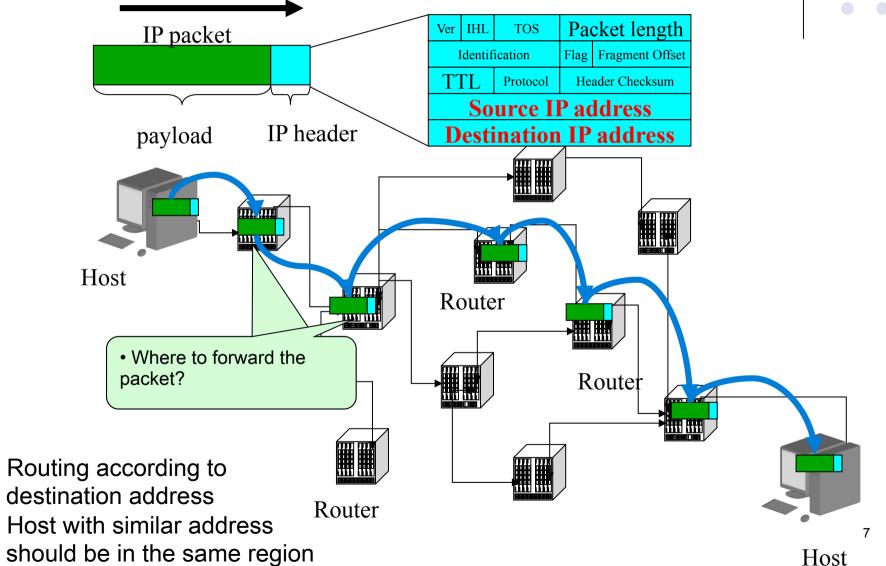




- Elements
 - host = end system;
 - subnetwork = a collection of hosts that are connected by layer-2 devices
 - Host of the same subnetwork have similar addresses: a common prefix
 - Routers: intermediate nodes inteconnect subnetworks:
- Packet forwarding
 - direct: inside a subnetwork hosts communicate directly without routers, layer-2 device (switch) delivers packets to hosts
 - indirect: between subnetworks one or several routers forward packets based on
 - structured address space
 - routing tables: aggregation of entries

IP Routing and forwarding







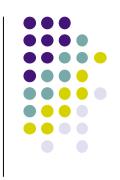


- Routing table
 - Each router have a routing table telling where to forward a packet

Network	Next-hop
10.0.0.0/24	Α
172.16.0.0/24	С
192.168.0.0/24	Direct

- Rule for sending packets (hosts, routers)
 - if the destination IP address has the same prefix as one of my interfaces, send directly to that interface
 - otherwise send to a router as given by the IP routing table





- Not reliable / fast
 - Sending data in "best effort" manner
 - No mechanism to recover error data at the receiver
 - When necessary, leave the upper layer (TCP) to ensure the data reliability.
- Packets are processed independently one of the other.

IP address

IP address classes

CIDR – Classless Inter-Domain routing

Subnet and netmask

Special IP addresses

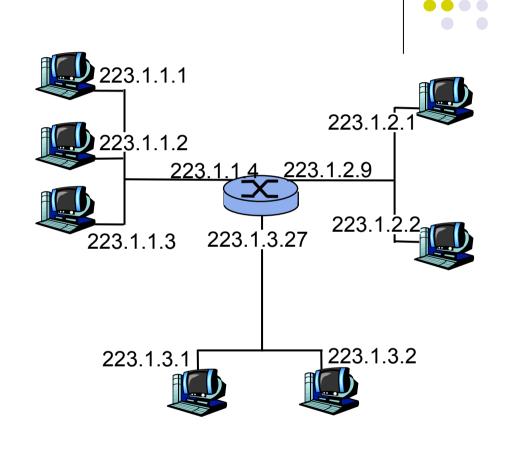


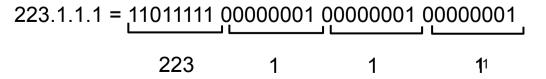
IP address (IPv4)

 IP address: A 32-bit number identifying uniquely a network interface

Interface:

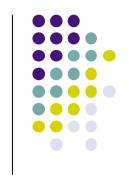
- router's typically have multiple interfaces
- host may have multiple interfaces
- IP addresses associated with interface, not host, router

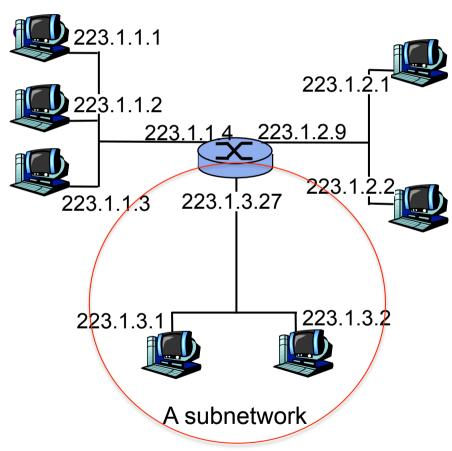


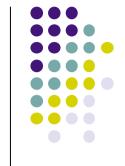




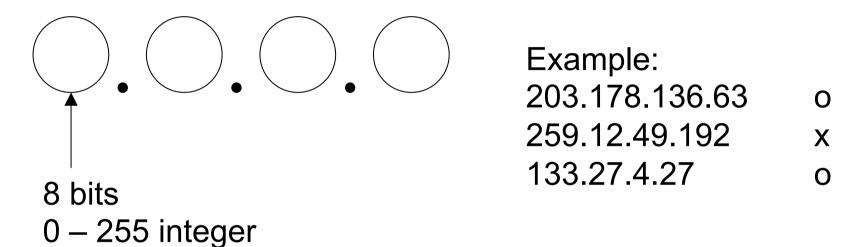
- For routing purpose, IP address of interfaces in the same subnetwork have the same prefix.
- What's a subnetwork?
 (from IP address perspective)
 - device interfaces with same prefix
 - can physically reach each other without intervening router (using layer 2 technology only)





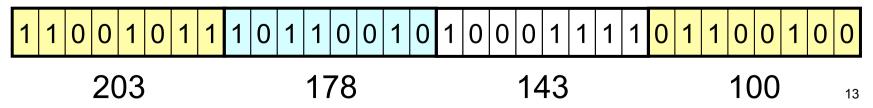


Dot notation



Use 4 x 8 bits describing a 32 bits address

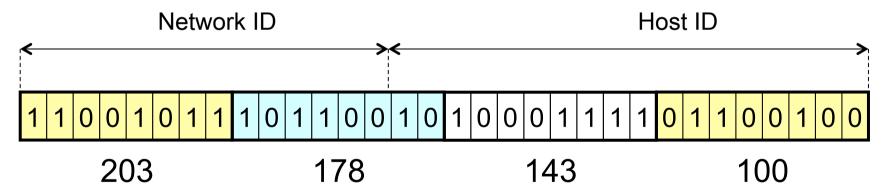
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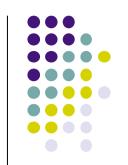


- IP address contains two parts
 - Host ID identify a host in a network
 - Network ID identify a network



- How to know which bits belong to network ID or host ID parts?
 - Use classful IP address
 - Use classless IP address
 CIDR

Classify IP addresses



	4	8bits			8bits	8bits	8bits	
C lass A	0		7 b it			Н	Н	Н
C lass B	1	0	6b it		N	Н	Н	
C lass C	1	1	0		5 b it	N	N	Н
C lass D	1	1	1	0		Мu	lticast	
C lass E	1	1	1	1	Res	serve fo	r future u	ıse

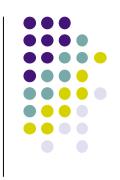
	# of network	# of hosts
C lass A	128	2^24
C lass B	16384	65536
C lass C	2^21	256





- Determine which classes these IP addresses belong to:
 - 10.10.10.9
 - 192.168.70.5
 - 129.60.4.7





- Inefficient use of addressing space
 - Hard classification of addressing space into classes (A, B, C, D, E) makes it is difficult to use all the address space

Solution...

- CIDR: Classless Inter Domain Routing
 - Network ID part will have variable length.
 - Length of Network ID part is specified in Network mask
 - Address notation: a.b.c.d/x, where x (mask) the number of bit of Network ID part.

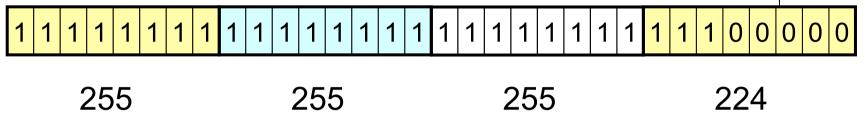




- Network mask divides the IP address into two parts
 - Part corresponding to Host ID
 - Part corresponding to Network ID
- IP addresses are assigned to hosts so that all hosts in the same network have the same Network ID part.
- Based on Network mask, it is possible to
 - Identify the network where an IP address belongs to
 - Calculate how many IP addresses available in the network associated with the mask.

Presentation of network mask





- 255.255.255.224
- /27
- 0xFFFFFe0

Last byte may be:

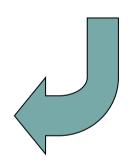


128 252

192 254

224 255

240



Calculation of network address

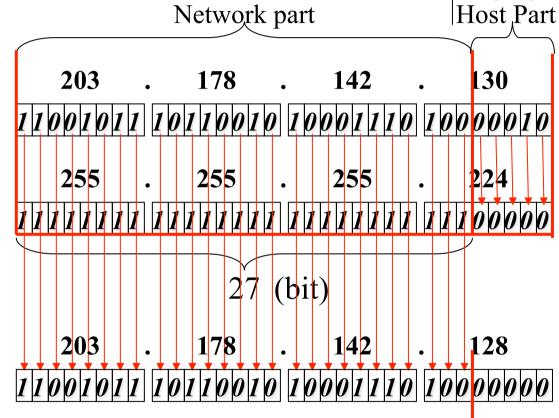


IP Address

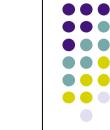
Netmask (/27)

AND

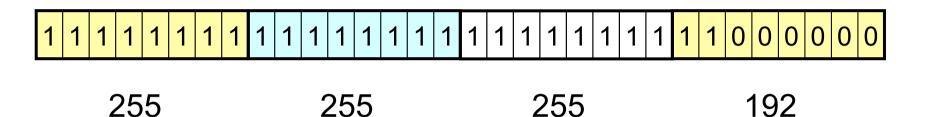
Network address



203.178.142.128/27



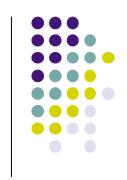
Calculation of network size

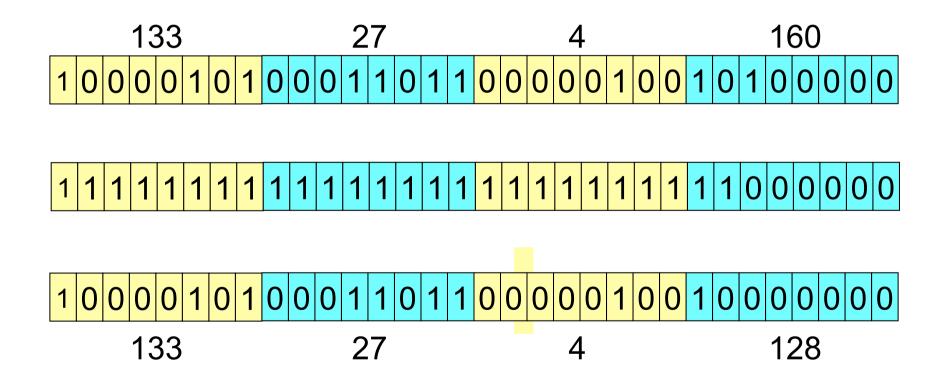


- Network size
 - Power of 2
- <u>RFC1878</u>

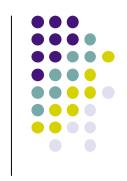
- In case of mask /26
 - Bits for Host ID = 6 bits
 - 2⁶=64 possible address:
 - 0 63
 - 64 127
 - 128 191
 - 192 255
 - Including network address and broadcast address

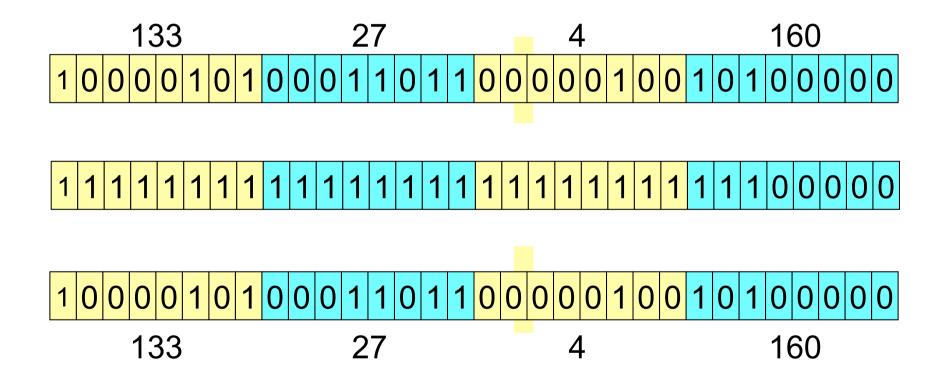
Network address or host address (1)





Network address or host address (2)





Different significations of IP address



- Network address
 - IP address assigned to a network
 - hostID contains all 0
- Host address
 - IP address assigned to a network card
- Broadcast address
 - Address used for sending data to all hosts in a network
 - All bit 1 in HostID part.

Exercice: IP address and network mask



- Which of the following IP addresses are host address, network address, broadcast address?
- (1) 203.178.142.128 /25
- (2) 203.178.142.128 /24
- (3) 203.178.142.127 /25
- (4) 203.178.142.127 /24
- Attn: With CIDR addressing, IP address should always coming with a network mask

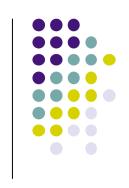


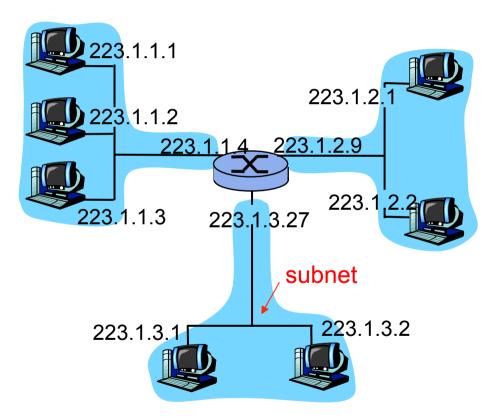
network

Subnet

Subnet is a part of a

- Hosts of a subnet communicate directly without reaching to layer 3.
- Usually is one department of an organization
- Design question: How to assign addresses of a network to subnets
 - Use a longer netmask





A network with 3 subnets.





```
11001000 00010111 00010000 00000000
          23.
200.
                     16.
                                        /24
11001000 00010111 00010000 00000000
200.
        23.
                    16.
                                       /25
11001000 00010111 00010000 10000000
200.
        23.
                    16.
                              128
                                       /25
```

Exercise: Dividing into subnets



- Given IP addresses in the range 200.23.16.0/24
- Need to organize into 4 subnets
 - Address of each subnetwork? Mask? Number of hosts/network
- General question: Need to create N subnets.
 Network address? Mask?
 - Each network contains 14 hosts /28
 - Each network contains 30 hosts _____ /27
 - Each network contains 31 hosts /26
 - Each network contains 70 hosts /25





- In theory
 - All between 0.0.0.0 ~ 255.255.255.255
 - Some special IP address (RFC1918)

	10.0.0/8		
Private address	172.16.0.0/12		
	192.168.0.0/16		
Loopback address	127.0.0.0		
Multipoet address	224.0.0.0		
Multicast address	~239.255.255.255		

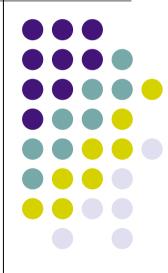
Self assigned IP address: 169.254.0.0/16





- Currently IPv4: 32 bits
 - 133.113.215.10 (IPv4)
- IPv6 is also widely used: 128bits
 - 2001:200:0:8803::53 (IPv6)
 - Fix 64 first bit for subnet ID, 64 last bit belongs to interface ID.
 - Security feature is integrated

IP package



Header of IP

total datagram length (words)

for



IP protocol version number

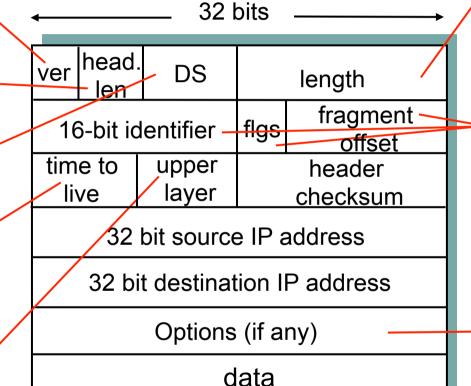
(bytes)

header length

QoS support

max number remaining hops (decremented at each router)

upper layer protocol to deliver payload to



(variable length,

typically a TCP

or UDP segment)

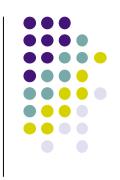
E.g. timestamp, record route taken, specify list of routers to visit.

fragmentation/

reassembly

IP header (1)

- Version (4 bits)
 - IPv4
 - IPv6
- Header length: 4bits
 - In word unit (4 bytes)
 - Min: 5
 - Max: 60







- DS (Differentiated Service : 8bits)
 - Old name: Type of Service
 - Used for QoS management by some router
 - Diffserv





- Length: total length including header (16 bits)
 - In bytes unit
 - Max: 65536
- 16 bits Identifier— ID of the packet
 - Used for identifying all fragments of the same packet when it is fragmented
 - Flag
 - Fragmentation offset offset of the first byte of the fragment in its original packet





- TTL, 8 bits Time to live
 - Maximum number of hops (router) the packet is allowed to travel
 - Max: 255
 - Router decreases TTL 1 unit when processing a packet
 - The packet will be destroyed when TTL reaches to 0
- Protocol upper layer protocol
 - Transport protocol (TCP, UDP,...)
 - Other network layer protocols that are encapsulated in IP packet (ICMP, IGMP, OSPF)

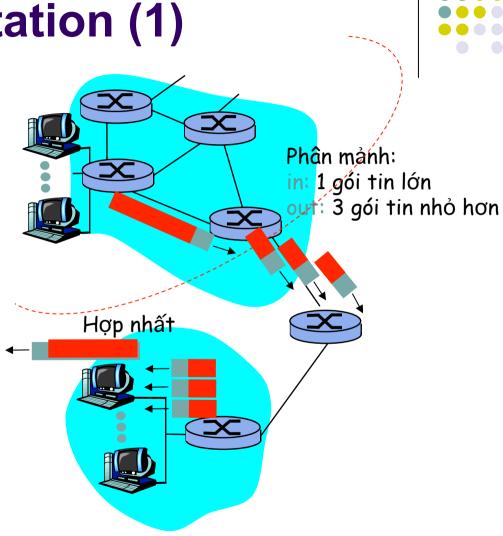




- Checksum: to detect corruption in the header of IPv4 data packets
- Source IP address
 - 32 bit, address of the sender
- Destination IP address
 - 32 bit, address of the receiver.

Packet fragmentation (1)

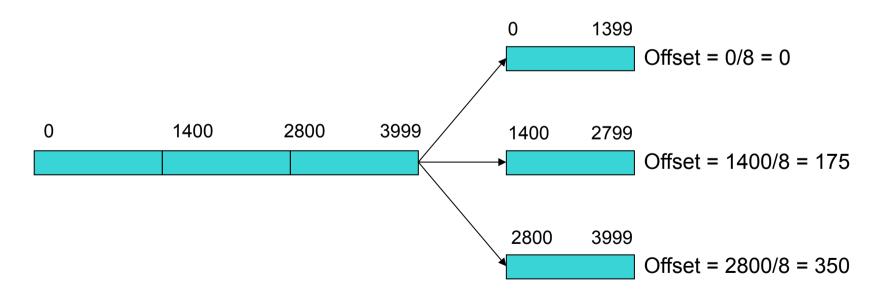
- Each link has a fixed MTU (Maximum transfering unit)
- Different media have different MTU
- If IP packet > MTU, it should be
 - Divided into small fragments
 - Gathered at the destination







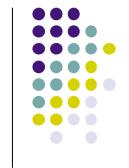
- Offset
 - Position of the fragment in the original packet
 - In 8 bytes units



Internet Control Message Protocol

Packet format Ping and Traceroute





Idea of ICMP (1)

- IP is unreliable, connectionless
 - Lack of supporting and error control mechanism
- ICMP is used in network layer for providing information exchange between sender and receivers
 - Error information: inform that a packet cannot reach a host, a network or a port.

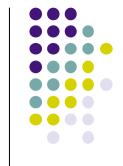




- Also in network layer but is "above" IP
 - ICMP message is encapsulated in IP
- ICMP message: Type, Code, with 8 first bytes of the error IP message

ICMP message

IP header ICMP message



IP header and Protocol field

Ver	HLEN	DS	Total Length			
Identification			Flags	Fragmentation offset		
TTL		Protocol	Header Checksum			
Source IP address						
Destination IP address						
Option						

Protocol:

1: ICMP

2: IGMP

6: TCP

17: UDP

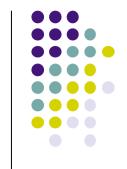
89: OSPF





- Type: type of ICMP message
- Code: cause of error
- Checksum
- Rest of header varies according on type

0	78 1516	31				
Туре	Code	Checksum				
Rest of the header						
Data						



Some ICMP message types

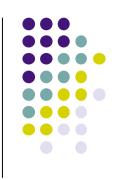
ICMP Message Type	Error-reporting messages	3	Destination Unreachable	
		4	Source quench (nguồn giảm tốc độ)	
		5	Redirection	
		11	Time exceeded	
		12	Parameter problem	
	Query messages	8 or 0	Echo reply or request	
		13 or 14	Time stamp request or reply	
		17 or 18	Address mask request or reply	
		9 or 10	Router advertisement or solicitation	





- ICMP always works transparently for users
- Users can use ICMP by using some debuging tools
 - ping
 - traceroute





- ping
 - Test a connection
 - Sender sends packet "ICMP echo request"
 - Receiver responses with "ICMP echo reply"
- Data field contains the time stamp when the packet is sent
 - For calculating RTT (round-trip time)





C:\Documents and Settings\hongson>ping www.yahoo.co.uk

Pinging www.euro.yahoo-eu1.akadns.net [217.12.3.11] with 32 bytes of data:

Reply from 217.12.3.11: bytes=32 time=600ms TTL=237

Reply from 217.12.3.11: bytes=32 time=564ms TTL=237

Reply from 217.12.3.11: bytes=32 time=529ms TTL=237

Reply from 217.12.3.11: bytes=32 time=534ms TTL=237

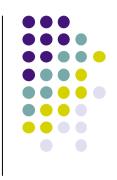
Ping statistics for 217.12.3.11:

Packets: Sent = 4, Received = 4, Lost = 0 (0% loss),

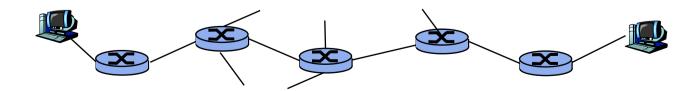
Approximate round trip times in milli-seconds:

Minimum = 529ms, Maximum = 600ms, Average = 556ms

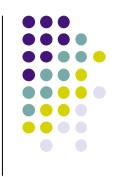




- Sender send many packets to receiver
 - First packet has TTL =1
 - Second packet has TTL=2, ...
- When packet number n arrives to nth router:
 - Router destroys the packer
 - Router send back an ICMP packet (type 11, code 0) containing IP address of the router
- Based on the reply message, the sender can calculate RTT

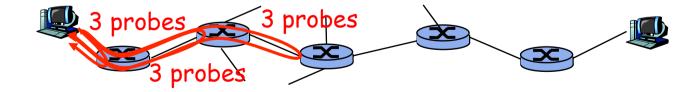






Termination condition

- When ICMP echo packet arrive to the destination
- When source receives ICMP "host unreachable" (type 3, code 3)





Traceroute: Example

C:\Documents and Settings\hongson>tracert www.jaist.ac.jp

Tracing route to www.jaist.ac.jp [150.65.5.208] over a maximum of 30 hops:

```
1 1 ms <1 ms 192.168.1.1
2 15 ms 14 ms 13 ms 210.245.0.42
3 13 ms 13 ms 13 ms 210.245.0.97
4 14 ms 13 ms 14 ms 210.245.1.1
5 207 ms 230 ms 94 ms pos8-2.br01.hkg04.pccwbtn.net [63.218.115.45]
6 * 403 ms 393 ms 0.so-0-1-0.XT1.SCL2.ALTER.NET [152.63.57.50]
7 338 ms 393 ms 370 ms 0.so-7-0-0.XL1.SJC1.ALTER.NET [152.63.55.106]
8 402 ms 404 ms 329 ms POS1-0.XR1.SJC1.ALTER.NET [152.63.55.113]
9 272 ms 288 ms 310 ms 193.ATM7-0.GW3.SJC1.ALTER.NET [152.63.49.29]
10 205 ms 206 ms 204 ms wide-mae-gw.customer.alter.net [157.130.206.42]
11 427 ms 403 ms 370 ms ve-13.foundry2.otemachi.wide.ad.jp [192.50.36.62]
12 395 ms 399 ms 417 ms ve-4.foundry3.nezu.wide.ad.jp [203.178.138.244]
13 355 ms 356 ms 378 ms ve-3705.cisco2.komatsu.wide.ad.jp [203.178.136.193]
14 388 ms 398 ms 414 ms c76.jaist.ac.jp [203.178.138.174]
15 438 ms 377 ms 435 ms www.jaist.ac.jp [150.65.5.208]
```

Trace complete.